# Survey of Degree Spectra of $High_n$ and $Non-low_n$ Degrees

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#### Spectrum of a Structure

**Defns**: For a countable structure S with domain  $\omega$ , the *Turing degree of* S is the Turing degree of S the the three the atomic diagram of S. The spectrum of S is

$$\operatorname{Spec}(\mathcal{S}) = \{ \operatorname{deg}(\mathcal{A}) : \mathcal{A} \cong \mathcal{S} \}$$

of all Turing degrees of copies of  $\mathcal{S}$ .

For a relation R on a computable structure  $\mathcal{M}$ , the spectrum of R,  $\mathrm{DgSp}_{\mathcal{M}}(R)$ , is

 $\{\deg(f(R)): f: \mathcal{M} \cong \mathcal{N} \& \mathcal{N} \text{ is computable}\}.$ 

#### Non-low Degrees

**Theorem** (Slaman; Wehner; Hirschfeldt): There exists a structure whose spectrum contains every Turing degree > 0, but not the degree 0.

This also holds with an arbitrary d in place of 0.

**Theorem** (GHKMMS): For each  $n \in \omega$ , there exists a structure whose spectrum contains precisely the non-low<sub>n</sub> degrees. Indeed, for each computable successor ordinal  $\alpha$ , there exists a structure with spectrum

 $\{ \deg(X) : (\exists \mathbf{d} \notin \mathbf{\Delta}^{\mathbf{0}}_{\alpha}) [\mathbf{d} \text{ is } \Delta^{\mathbf{0}}_{\alpha} \text{ relative to } \mathbf{X}] \}.$ 

### High Degrees

Given a structure  $\mathcal{A}$ , the technique of (GHKMMS) builds, for each successor ordinal  $\alpha$ , a structure  $\mathcal{B}$  such that

 $\boldsymbol{c} \in \operatorname{Spec}(\mathcal{B}) \iff \boldsymbol{c}^{(\alpha)} \in \operatorname{Spec}(\mathcal{A}).$ 

**Fact**: For every d, there is a structure  $\mathcal{A}_d$  with spectrum  $\{c : c \geq_T d\}$ .

So with  $\alpha = 1$  and  $\mathbf{d} = \mathbf{0}''$ , this gives a structure  $\mathcal{B}$  whose spectrum contains exactly the high-or-above degrees (those  $\mathbf{c}$  with  $\mathbf{c}' \geq_T \mathbf{0}''$ ). Likewise for high<sub>n</sub>, and even high<sub> $\alpha$ </sub> (with  $\alpha \notin \mathbf{LOR}$ ). This extends a known result...

#### Spectrum of high degrees

**Proposition** (Harizanov-Miller): There exists a relation R on a computable dense linear order  $\mathbb{Q}$  with

 $\mathrm{DgSp}_{\mathbb{Q}}(R) = \{ \mathbf{d} : \mathbf{d}' \geq_{\mathbf{T}} \mathbf{0}'' \}.$ 

**Corollary**: There exists a structure with this same spectrum.

**Corollary**: Not all spectra of unary relations on  $(\mathbb{Q}, <)$  can be realized as spectra of linear orders. **Proof**: By a theorem of Knight (1986), **0'** is the only possible jump degree of a linear order.

#### How About Linear Orders?

For Boolean algebras, the spectrum  $\{\boldsymbol{d}:\boldsymbol{d}>\boldsymbol{0}\}$  is impossible.

- Jockusch-Soare: For every c.e. d > 0, there exists a linear order whose spectrum contains d but not 0.
- Downey/Sectapun: Extension to any d with  $0 < d \le 0'$ .
- Knight: Extension to any d > 0.
- M.: There is a single linear order whose spectrum contains all d with  $0 < d \le 0'$ , but not 0.
- Frolov: For each  $n \in \omega$ , there is a linear order whose spectrum contains all non-low<sub>n</sub> degrees  $\leq \mathbf{0}'$  but no low<sub>n</sub> degrees.

**Question**: Can a linear order have a spectrum of precisely the non-low<sub>n</sub> degrees?

#### Where Next?

Frolov's result builds an order  $\mathcal{L}$  by relativizing Miller's result to  $\mathbf{0}^{(n)}$ , so that  $\operatorname{Spec}(\mathcal{L})$  contains all d with  $\mathbf{0}^{(n)} < d \leq \mathbf{0}^{(n+1)}$ , but not  $\mathbf{0}^{(n)}$ . Then it applies the relativization of a theorem of Downey and Knight: A linear order  $\mathcal{L}$  is  $\Delta_2^0$  iff  $(\eta + 2 + \eta) \cdot \mathcal{L}$  is computable. So the order  $\mathcal{L}_n = (\eta + 2 + \eta)^n \cdot \mathcal{L}$  has all non-low<sub>n</sub> degrees below  $\mathbf{0}'$  in its spectrum, but no low<sub>n</sub> degree.

#### Spectral Universality

An embedding  $f : \mathcal{A} \hookrightarrow \mathcal{B}$  preserves the spectrum if  $\operatorname{Spec}(A) = \operatorname{DgSp}_{\mathcal{B}}(f(\mathcal{A})).$ 

A computable model  $\mathcal{B}$  of a theory T is spectrally universal for T if every countable model  $\mathcal{A}$  of Tembeds into  $\mathcal{B}$  via some f preserving the spectrum.

Many (but not all!) computable Fraïssé limits of theories turn out to be spectrally universal for those theories. Examples:

- Countable dense linear order.
- Random graph.
- Countable atomless Boolean algebra.

Counterexample:

• Algebraic closure of the field  $\mathbb{Z}/(p)$ .

#### Structures vs. Relations

**Corollary** (of the spectral universality of the random graph): The spectra of countable graphs are precisely the spectra of unary relations on the random graph.

We saw above that this does *not* hold of linear orders. Spectral universality of the countable DLO shows that every spectrum of a LO is the spectrum of a unary relation on the DLO, but the converse is false.

#### New Possible Counterexample

**Theorem** (M.): There exists a unary relation  $\tilde{R}$ on the countable DLO ( $\mathbb{Q}, \prec$ ) whose degree spectrum contains the non-low degrees:

$$\mathrm{DgSp}_{\mathbb{Q}}(\tilde{R}) = \{ \mathbf{d} : \mathbf{d'} >_{\mathbf{T}} \mathbf{0'} \}.$$

For a given finite set  $F = \{n_1, n_2, \dots, n_k\} \subset \omega$ and  $a \prec b$  in  $\mathbb{Q}$ , define  $\tilde{R}$  on (a, b) for F as follows:

Wehner's construction gives a family  $\mathcal{F}$  of finite sets F, uniformly enumerable by any degree  $> \mathbf{0}'$ , but not by  $\mathbf{0}'$ .

For each  $F \in \mathcal{F}$ , in each order, define  $\tilde{R}$  as above for this F on densely many intervals (a, b) in  $\mathbb{Q}$ .

#### Does this work?

This  $\tilde{R}$  on  $(\mathbb{Q}, \prec)$  gives a potential second counterexample to the converse of spectral universality for linear orders.

Question: Does there exist a linear order with spectrum  $\{c : c' >_T 0'\}$ ?

Notice that the restriction of  $\prec$  to  $\tilde{R}$  does *not* yield such an order.

**Example:** Fix  $r_0 = e$  and  $r_{i+1} = e^{r_i}$ . Given  $S \subseteq \omega$ , let  $F_S$  be the closure of  $\mathbb{Q}(r_t \mid t \in S)$  under square roots of positive elements. We claim that

$$\operatorname{Spec}(F_S) = \{ \boldsymbol{d} : S \text{ is } \Sigma_2^0 \text{ in } \boldsymbol{d} \}.$$

**Cor.**: For any  $A \subseteq \omega$ , there is a field whose spectrum contains precisely those d with  $A \leq d'$ .

$$\operatorname{Spec}(F_{A'}) = \{ \boldsymbol{d} : (\exists D \in \boldsymbol{d}) A' \leq_1 D'' \}$$
$$= \{ \boldsymbol{d} : (\exists D \in \boldsymbol{d}) A \leq_T D' \}$$

As a relation on its algebraic closure,  $F_S$  has the same spectrum  $\{\boldsymbol{d}: S \text{ is } \Sigma_2^0 \text{ in } \boldsymbol{d}\}.$ 

So the high degrees form the spectrum of a field, and also the spectrum of a subfield of the algebraic closure.

## $\operatorname{\mathbf{Spec}}(F_S) = \{ \boldsymbol{d} : S \text{ is } \Sigma_2^0 \text{ in } \boldsymbol{d} \}.$

 $\subseteq$ : If  $E \cong F_S$ , then S is the set

 $\{t \in \omega : (\exists x \in E) (\forall q \in \mathbb{Q}) [q < r_t \leftrightarrow q \prec x \text{ in } E]\}.$ 

The order  $\prec$  on E is E-computable, by the closure of E under square roots of positive elements.

 $\supseteq$ : If  $S \leq_1 \operatorname{Fin}^D$ , let  $t \in S$  iff  $W_{h(t)}^D$  is finite. Start building  $F_{\omega}$  (the field containing all  $r_t$ ). Each time  $W_{h(t)}^D$  receives an element, make the old  $r_t$ become rational and add a new  $r_t$  to replace it.