On the question of consistence of the semantic μ -prediction

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Current decisions are made in two-valued classical logic, so consistency of probabilistic theories/predictions (statistical ambiguity problem) is a very important question of AI.

Note that any ϕ should be examined both with its negation: each of them may be specific in prediction of some ψ , e.g. $\mu(\psi | \phi) > \mu(\psi | \neg \phi)$ or $\mu(\psi | \phi) < \mu(\psi | \neg \phi)$, where $\mu(\phi) > \mu(\neg \phi)$, for instance.

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Let \mathfrak{L} be a first-order language of a finite signature.

Allow literals (not only atoms) to appear in a classical logic programming structures of rule, fact and query; denote corresponding sets as Rule_£, Fact_£ and Query_£.

definition

Language

Binary relation $C_1 \succ C_2$ (read " C_1 is more general than C_2 ") between $C_1 \equiv (A_1 \Leftarrow B_1, ..., B_n), C_2 \equiv (A_2 \Leftarrow D_1, ..., D_m)$ in Rule_£ takes place iff there exist a substitution θ such that $\{B_1\theta, ..., B_n\theta\} \subseteq \{D_1, ..., D_m\}, A_1\theta \equiv A_2$ and $\not\vdash C_1 \equiv C_2$.

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Probability over ground sentences

• Let \mathfrak{G}^* be a countable class of observed first-order structures appearing in practice; $\mathfrak{G} \subset \mathfrak{G}^*$ is a general sampling consisting of well-studied models.

• Being given \mathfrak{G} we compute a probability measure P over \mathfrak{G}^* with some trusting interval value $\varepsilon > 0$ (according to Kolmogorov); here mathematical statistics is applied.

• Assume $\mu(\phi) = \mathsf{P}(\{\mathfrak{A} | \mathfrak{A} \models \phi\})$, where ϕ is a closed formulae.

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Let $\Theta^{\rm o}$ be a set of all ground substitutions.

Probability

Introduction

Consistence

Probability of a ground instance of rule is defined as conditional $\mu (A \Leftarrow B_1 \land ... \land B_n) = \mu (A | B_1 \land ... \land B_n) = \frac{\mu (A \land B_1 \land ... \land B_n)}{\mu (B_1 \land ... \land B_n)}$

 $\mathtt{Rule}^{\mu}_{\mathfrak{L}} \rightleftharpoons \{ C \mid \text{for some } \theta \in \Theta^{\mathrm{o}} \text{ probability of } C\theta \text{ is determined} \}$

$$\underline{\mu}(\mathbf{C}) \rightleftharpoons \inf \left\{ \mu\left(\mathbf{C}\theta\right) \mid \theta \in \Theta^{\mathrm{o}} \text{ and } \mathbf{C}\theta \in \mathtt{Rule}_{\mathrm{L}}^{\mu} \right\},$$
where $\mathbf{C} \in \mathtt{Rule}_{\mathfrak{L}}^{\mu}$

 $Fact_o$ is a set of ground atoms allowing verification in any $\mathfrak{B} \in \mathfrak{G}^*$; a complete set of alternatives is

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$$\mathtt{Fact}_{\mathrm{o}}^* = \mathtt{Fact}_{\mathrm{o}} \cup \{ \neg \mathrm{A} \, | \, \mathrm{A} \in \mathtt{Fact}_{\mathrm{o}} \}$$

definition (E.E. Vityaev, S.O. Smerdov)

A rule $C \equiv (A \Leftarrow B_1 \land ... \land B_n)$ is called *the best prediction rule for some literal* H iff the following conditions are hold:

- $$\begin{split} &i. \text{ there exist } \theta \in \Theta^{\mathrm{o}} \text{ such that } \mathrm{A}\theta \equiv \mathrm{H}\theta \text{, } \{\mathrm{B}_{1}\theta,...,\mathrm{B}_{n}\theta\} \subseteq \mathtt{Fact}_{\mathrm{o}}^{*}\text{,} \\ &\mu\left((\mathrm{B}_{1}\wedge...\wedge\mathrm{B}_{n})\,\theta\right) \neq \texttt{0} \text{ and } \underline{\mu}\left(\mathrm{C}\right) > \mu\left(\mathrm{H}\theta\right)\text{;} \end{split}$$
- *ii.* maximum of $\underline{\mu}(\cdot)$ is achieved on C among rules satisfying (i);
- *iii.* it can't be generalized without loosing (i-ii).

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• We denote by $Prdct_{\mathfrak{L}}^{\mu,0}$ the obtained set of described ground instances (over all literals H).

• Data(\mathfrak{B}) is a set of actual facts for 1-st order model $\mathfrak{B} \in \mathfrak{G}^*$, i.e. consistent subset of Fact^{*}_o (not necessary maximal).

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A set of literals S is called μ -concurred iff $P(\{\mathfrak{A} \mid \mathfrak{A} \models S\}) \neq 0$.

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theorem

Let some ground atom H be semantically μ -predicted by ground instance $C_{pos} \in \text{Prdct}_{\mathfrak{L}}^{\mu,0}$ of the best rule C_1 ($C_{pos} \equiv C_1\theta_{pos}$), while $\neg H$ is predicted by $C_{neg} \in \text{Prdct}_{\mathfrak{L}}^{\mu,0}$ ($C_{neg} \equiv C_2\theta_{neg}$). Then the set of atoms from premises of C_{pos} and C_{neg} is not μ -concurred.

Denote by $\Gamma_{\mathfrak{B}}$ the following set of rules and data $\left\{ \operatorname{B}_{1} \wedge ... \wedge \operatorname{B}_{n}
ightarrow \operatorname{A} | \operatorname{A} \Leftarrow \operatorname{B}_{1} \wedge ... \wedge \operatorname{B}_{n} \in \operatorname{Prdct}_{\operatorname{L}}^{\mu,0} \right\} \cup \operatorname{Data}(\mathfrak{B})$

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Let $Data(\mathfrak{B})$ be μ -concurred. Then minimal theory containing $\Gamma_{\mathfrak{B}}$ is logically consistent.

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Consistence	Theorems

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