Ideals in the Turing degrees Examples via randomness; upper bounds

> André Nies The University of Auckland

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For instance, each proper Σ_3^0 ideal has a low₂ upper bound.

Part I

Background on ideals

The ideal lattice of an usl U

- Let $(U, \leq \vee)$ be an uppersemilattice (usl).
- A set *I* ⊆ *U* is an ideal if *I* is closed downwards and under the join operation ∨.
- An **upper bound** of an ideal *I* is a degree **b** such that $I \subseteq [0, b]$.

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Some Facts:

- The set of ideals of *U* is a lattice, where the meet of *I*, *J* is the intersection, and the join of *I*, *J* is the ideal generated by *I* ∪ *J*.
- An ideal *I* is called **proper** if $I \neq U$.
- Each *u* ∈ *U* determines the ideal {*x*: *x* ≤ *u*}, called a principal ideal.

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- to study quotient structures.
- There are many examples, because several algebraic operators in usl turn sets into ideals.
- some important classes are ideals, such as "cappable" in the c.e. degrees, "*K*-trivial" in the Δ⁰₂, and the c.e. degrees.
- Ideals form an abstract framework for some lowness properties.

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If U has a largest element 1, then the core of $U - \{1\}$ is

 $\{x: \forall d < 1 [x \lor d < 1]\}$ = non-cuppable.

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- Applying this, Yang Yue and Yu Liang found a few more examples of definable ideals: for instance, the ideal generated by the non-bounding degrees.

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- Each principal ideal is Σ_4^0 .
- For $k \ge 4$, the $\sum_{k=1}^{0} \frac{1}{2} \sum_{k=1}^{n} \frac{1}{2} \sum_{k=$

Classes of ideals in the c.e. degrees

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It is not hard to show that the converse implications fail:

- Let a < 1 be a non-low₂ c.e. degree. Then [0, a] is u.g. but not Σ₃⁰.
- If b ≠ 0, then the principal ideal [0, b] has a maximal subideal I that is Δ⁰₄(b). Now choose b low. Then I is Σ⁰₄ but not u.g. as we'll see later.

Part II

Ideals via randomness

Strongly jump traceable sets

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- The set *A* is called strongly jump traceable if for each order function *h*, there is a c.e. trace (*T_x*)_{x∈N} with bound *h* such that, whenever *J^A(x)* it is defined, we have

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Definition of cost functions

Definition A cost function is a computable function

$$c: \mathbb{N} \times \mathbb{N} \to \{x \in \mathbb{Q}: x \ge 0\}.$$

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When building a computable approximation of a $\Delta_2^0 \text{ set } A$, we view c(x, s) as the cost of changing A(x) at stage s.

Obeying a cost function

We want to make the **total** cost of changes, taken over all *x*, **finite**.

Definition

The computable approximation $(A_s)_{s \in \mathbb{N}}$ obeys a cost function *c* if

 ∞ > $\sum_{x,s} c(x,s) [x < s \& x \text{ is least s.t. } A_{s-1}(x) ≠ A_s(x)].$

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We write Models(c) for the c.e. sets *A* that obey *c*. For monotonic *c*, this class is closed under \oplus .

Basic existence theorem

We say that a cost function c satisfies the limit condition if

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Theorem (Kučera, Terwijn 1999; D,H,N,S 2003; ...)

If a cost function *c* satisfies the limit condition, then some simple set *A* obeys *c*.

The ideal $\mathcal{I}(Y)$

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- $\mathcal{I}(Y)$ induces an ideal in the c.e. degrees.
- By Kučera's Theorem, if the Δ⁰₂ set Y is ML-random then I(Y) contains a promptly simple set.
- [Greenberg, N.] For each Δ⁰₂ set Y there is a cost function c_Y with the limit condition such that

 $A \models c_Y \& Y$ ML-random $\Rightarrow A \leq_T Y$.

That is, $Models(c_Y) \subseteq \mathcal{I}(Y)$ for ML-random Y.

Basis Theorems

Recall: $\mathcal{I}(Y) = \{A \text{ c.e. } : A \leq_T Y\};$

Models(c) is the class of c.e. sets A such that A obeys c.

Theorem

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Let \mathcal{P} be a non-empty Π_1^0 class (such as a class of ML-randoms). Let c be a monotonic cost function with the limit condition. (i) [N.] There is a Δ_2^0 set $Y \in \mathcal{P}$ such that $Models(c) \not\subseteq \mathcal{I}(Y)$. (ii) [Greenberg, Hirschfeldt, N] There is a Δ_2^0 set $Z \in \mathcal{P}$ such that $\mathcal{I}(Z) \subseteq Models(c)$.

In (i) one builds $Y \in \mathcal{P}$ and a c.e. set $A \models c$ such that $A \not\leq_T Y$. (ii) says that for each c.e. set $A \leq_T Z$ we have $A \models c$.

Diamond Classes

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2[№] denotes Cantor space with the uniform (coin-flip) measure. We define ideals in the c.e. degrees as the lower bounds of classes of ML-random sets.

For a null class $\mathcal{H} \subseteq \mathbf{2}^{\mathbb{N}}$, we let

 \mathcal{H}^{\Diamond} = the c.e. sets Turing below each ML-random set in \mathcal{H} .







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- (Hirshfledt/Miller) For each null Σ⁰₃ class *H*, there is a promptly simple set in *H*[◊].
- In the interesting case that there is a ML-random set Y ≥_T Ø' in H, we have H[◊] ⊆ base for ML-random (= K-trivial).

Lowness, Highness

For a set X, we let X' denote the halting problem relative to X.

- Recall that $Z \subseteq \mathbb{N}$ is low if $Z' \leq_T \emptyset'$, and Z is high if $\emptyset'' \leq_T Z'$.
- These classes are "too big" in this context: we have

 $(low)^{\diamond} = (high)^{\diamond} = computable.$

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(For instance, $(high)^{\diamond}$ = computable because there is a minimal pair of high ML-random sets.)

 So we will try somewhat smaller classes, replacing ≤_T by the stronger truth-table reducibility ≤_{tt}. Diamond classes coinciding with SJT_{c.e.}

Definition (Mohrherr 1986)

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Theorem (Greenberg, Hirschfeldt and Nies, to appear)

A c.e. set A is strongly jump traceable

- ↔ A is Turing below each superlow ML-random set
- ↔ A is Turing below each superhigh ML-random set .

Diagram: $SJT_{c.e.}$ means computed by many oracles



Remember that in an usl U, the **core** of $S \subseteq U$ is

 $\{x \in U \colon \forall d \in S [x \lor d \in S]\}.$

As a corollary of $SJT_{c.e.} \subseteq (superlow)^{\diamond}$, we have that (at least on the c.e. sets), SJT is contained in the core of the superlow sets.

Theorem (Greenberg and Nies (2008))

Suppose the c.e. set *A* is strongly jump traceable. Then (*) $\forall X$ superlow $[X \oplus A \text{ is superlow}]$.

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Is (*) a characterization of SJT_{c.e.}? Is the ideal induced by (*) at least contained in the K-trivials?

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If we restrict (*) to c.e. sets X, then it properly contains $SJT_{c.e.}$

(Diamondstone and No. to appear)

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Open questions on ideals between $SJT_{c.e.}$ and *K*-trivial

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- Here AED is the class of almost everywhere dominating sets *D* of Dobrinen and Simpson: for almost all sets *X*, each function *f* ≤_T *X* is dominated by a function *g* ≤_T *D*.

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- Here AED is the class of almost everywhere dominating sets *D* of Dobrinen and Simpson: for almost all sets *X*, each function *f* ≤_T *X* is dominated by a function *g* ≤_T *D*.
- For the highness properties, there are proper implications

Turing-complete \Rightarrow AED \Rightarrow superhigh.

$(AED)^{\Diamond}$ properly contains $SJT_{c.e.}$

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- However, $(AED)^{\diamond}$ may coincide with *K*-trivial.
- This would imply that the classes **ML-coverable** and **ML-noncuppable** also coincide with *K*-trivial.

Classes of c.e. sets between SJT_{c.e.} and K-trivial



(The dashed arrows may be coincidences.)

- A is ML-coverable if $A \leq_T Y$ for some ML-random $Y \geq_T \emptyset'$.
- A is ML-noncuppable if

 $\emptyset' \leq_T A \oplus Y$ for ML-random Y implies $\emptyset' \leq_T Y$.

Inside SJT_{c.e.}

Work in progress with Diamondstone and Hirschfeldt shows: The class

 $(\omega^{\omega}$ -c.e.)

is a nontrivial proper subclass of $SJT_{c.e.}$.

BREAK

Part III

Upper bounds for ideals (joint with G. Barmpalias)

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What can we say about upper bounds of / in the c.e. degrees? Motivation: often / is a lowness property. In this case we would expect results on upper bounds.

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More on the leading question

- By the Thickness Lemma every proper u.g. ideal has an incomplete upper bound.
- What can we say about upper bounds of a proper Σ⁰₃ ideal?
- The Π_4^0 ideal of cappable degrees has no incomplete upper bound.
- How about bounds for a proper Σ⁰₄ ideal?

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- We have a Ø" construction with a tree of strategies to read a low₂-ness index for the upper bound off the true path.

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- if $n \in \operatorname{Tot}^{Y_k}$ then $\operatorname{deg}(U_{k,n}) \in \mathbb{I}$
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This is done by attempting to enumerate \emptyset' into the $U_{k,n}$. At stage *s*, for each n, k < s:

if $v \in \emptyset'_s$ and $\Phi_n^{Y_k}(v) \uparrow [s]$, enumerate v into $U_{k,n}$.

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In order to build an incomplete upper bound, we build *B* meeting the requirements

$$C_{\langle e,n\rangle}: V_{e,n}^H = \mathbb{N} \Rightarrow W_e \leq_T B \oplus H.$$

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We make **B** Turing incomplete, by meeting the requirements

 $N_m: \quad \emptyset' = \Phi^B_m \Rightarrow \exists k \exists e_0, \dots, e_{k-1} [\emptyset' \leq_T \oplus_{i < k} W_{e_i} \oplus H \quad \& \forall i \ W_{e_i} \in \mathcal{I}].$

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This condition says that, if *B* is complete, then the ideal given by \mathcal{I} is not proper. The sets W_{e_i} , i < k, will be the members of \mathcal{I} that are coded into *B* through higher priority requirements.

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Ideals in the Turing degrees

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Corollary No proper Σ_4^0 ideal is prime.

For, pick an incomplete upper bound of the ideal. Welch 1981 shows that there is a minimal pair of degree none of which are below this upper bound.

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- This uses the Uniform Low₂-ness Lemma combined with a Sacks Splitting type technique.
- We also see now that I above is not uniformly generated: else it would already be Σ⁰₃.

Some open questions on ideals

- Is every Σ₄⁰ ideal I the intersection of the principal ideals it is contained in? (This would strengthen our result that I has an incomplete upper bound.)
- For k ≥ 4, is the class of principal ideals definable in the lattice of Σ⁰_k ideals? Natural elementary differences for k ≥ 4?
- Let K be the ideal of K-trivial degrees. Are there c.e. degrees a, b such that K = [0, a] ∩ [0, b]?

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